# Colorful Graph Theory 

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## Sunway TaihuLight



## Sunway TaihuLight

- Ranked \#1 in the TOP500 list in March 2018 as the fastest supercomputer
- 93 petaflops $=93 \cdot 10^{15}$ flops (floating point operations per second)
■ 10,649,600 CPU cores


## Parallel processing

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- Natural solution: more processors

■ Parallel processing: computations executed at the same time

## Simple Example

- 8 ice hockey teams;

■ Each team plays each team;

- Every day one match per team;
- We have 7 days;


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## Basics

■ For a graph $G=(V, E)$, a $k$-edge-coloring is a function

$$
f: E \rightarrow\{1,2, \ldots, k\}
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■ Upper bound?
Theorem 1 (Vizing, 1964)
For every (simple) graph G

$$
\Delta(G) \leq \chi^{\prime}(G) \leq \Delta(G)+1
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## Bipartite and complete graphs

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- For complete graphs $K_{2 k}$, we have $2 k-1$ disjoint perfect matchings; we assign the same color to all edges of a matching, so:

$$
\chi^{\prime}\left(K_{2 k}\right)=\Delta\left(K_{2 k}\right)=2 k-1
$$

Complete graphs of odd order need additional color:

$$
\chi^{\prime}\left(K_{2 k+1}\right)=\Delta\left(K_{2 k+1}\right)+1=2 k+1 .
$$

## Adding assumptions

- How do the bounds for chromatic index change if we add additional assumptions to the coloring?


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- We will focus on three types:
- Acyclic edge-coloring;
- Strong edge-coloring;
- Locally irregular edge-coloring.

Acyclic edge-coloring

## Acyclic edge-coloring

- An acyclic $k$-edge-coloring of a graph $G$ is a proper $k$-edge-coloring where the edges of every cycle are assigned at least three distinct colors, i.e., there are no bichromatic cycles.


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- So, $\chi_{a}^{\prime}\left(K_{4}\right)=5=\Delta\left(K_{4}\right)+2$


## The Conjecture

Conjecture 3 (Fiamčík, 1978; Alon, Sudakov, Zaks, 2001)
For every graph G it holds

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- Conjecture 3 is not confirmed even for complete graphs!


## Perfect 1-factorization

## Conjecture 4 (Kotzig, 1964)

For every $n \geq 2, K_{2 n}$ can be decomposed into $2 n-1$ perfect matchings such that the union of any two matchings forms a hamiltonian cycle in $K_{2 n}$.

- Closely related to acyclic edge-colorings.
- If the Conjecture 4 is true, the removal of one vertex from $K_{2 n}$ results in an acyclic edge coloring of $K_{2 n-1}$ with $2 n-1=\Delta\left(K_{2 n-1}\right)+1$ colors, which is optimal.


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- If $K_{n+1}$ has perfect 1-factorization, then $K_{n, n}$ has it also.

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## Theorem 5 (Giotis et al., 2017)

For every graph G it holds

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## Theorem 6 (Alon, Sudakov, Zaks, 2001)

For every graph $G$ with girth at least $C \Delta(G) \log \Delta(G)$, for a constant $C$, it holds

$$
\chi_{a}^{\prime}(G) \leq \Delta(G)+2 .
$$

## Subcubic graphs

- The notion of acyclic colorings was first introduced in 1973 by Grűnbaum for the vertex version. In 1979, Burnstein proved that 5 colors suffice for acyclic vertex coloring of every graph $G$ with $\Delta(G) \leq 4$.


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- The maximum degree of the line graph $L(G)$ of a subcubic graph $G$ is at most 4...


## Corollary 7 (Burnstein, 1979)

Let $G$ be a subcubic graph. Then

$$
\chi_{a}^{\prime}(G) \leq 5 .
$$

## Planar graphs

■ Another candidate class of graphs to confirm Conjecture 4

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Theorem 8 (Wang, Zhang, 2017+)
Let $G$ be a planar graph. Then

$$
\chi_{a}^{\prime}(G) \leq \Delta(G)+6 .
$$

## Planar graphs - $\Delta$ colors

■ Cohen, Havet and Müller conjectured that every planar graph $G$ with large enough maximum degree has $\chi_{a}^{\prime}(G)=\Delta$ (note the analogy to Vizing's conjecture)

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## Theorem 9 (Cranston, 2017+)

Let $G$ be a planar graph with $\Delta(G) \geq 4.2 \cdot 10^{14}$. Then

$$
\chi_{a}^{\prime}(G)=\Delta(G)
$$

## Planar graphs - $\Delta$ colors

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## Theorem 10 (Húdak et al., 2012)

Let $G$ be a planar graph with girth $g$ and maximum degree $\Delta$. Then $\chi_{a}^{\prime}(G)=\Delta$ if one of the following conditions holds:

- $\Delta \geq 3$ and $g \geq 12$, or
- $\Delta \geq 4$ and $g \geq 8$, or
- $\Delta \geq 5$ and $g \geq 7$, or
- $\Delta \geq 6$ and $g \geq 6$, or
- $\Delta \geq 10$ and $g \geq 5$.


## Strong edge-coloring

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■ Distance between edges: distance between corresponding vertices in the line graph (adjacent edges are at distance 1)

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- The smallest $k$ for which $G$ admits a strong $k$-edge-coloring is the strong chromatic index of $G, \chi_{s}^{\prime}(G)$.
- Strong edge coloring of $G$ is a vertex 2-distance coloring of its line graph $L(G)$

$$
\chi_{s}^{\prime}(G)=\chi\left(L(G)^{2}\right)
$$

Example


Example


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$\chi_{s}^{\prime}(G)=7$

## The Conjecture

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## Conjecture 11 (Erdős, Nešetřil, 1985)

For every graph G it holds

$$
\chi_{s}^{\prime}(G) \leq\left\{\begin{array}{cl}
\frac{5}{4} \Delta(G)^{2}, & \Delta(G) \text { is even; } \\
\frac{1}{4}\left(5 \Delta(G)^{2}-2 \Delta(G)+1\right), & \Delta(G) \text { is odd }
\end{array}\right.
$$

## The Conjecture

- The bounds in Conjecture 11 are tight for every $\Delta$ :



## The Conjecture

The construction of graphs achieving the conjectured bound:

- For even $\Delta$ replace each vertex of a 5 -cycle with $\frac{\Delta}{2}$ vertices;

■ For odd $\Delta$ replace two consecutive vertices of a 5 -cycle with $\frac{\Delta+1}{2}$ vertices and the others with $\frac{\Delta-1}{2}$ vertices.

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## Theorem 12 (Bonamy, Perret, Postle, 2017+)

For every graph $G$ with sufficiently large maximum degree it holds

$$
\chi_{s}^{\prime}(G) \leq 1.835 \Delta(G)^{2}
$$

## Subcubic \& Subquartic graphs

## Theorem 13 (Andersen, 1992)

Let $G$ be a graph with $\Delta(G)=3$. Then,

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\chi_{s}^{\prime}(G) \leq 10 .
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Theorem 14 (Cranston, 2006)
Let $G$ be a graph with $\Delta(G)=4$. Then,

$$
\chi_{s}^{\prime}(G) \leq 22 .
$$

## Bipartite graphs

Conjecture 15 (Faudree et al., 1990)
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## Conjecture 16 (Brualdi, Quinn Massey, 1993)

If $G$ is bipartite graph with maximum degree of partite sets $\Delta_{1}$ and $\Delta_{2}$, then

$$
\chi_{s}^{\prime}(G) \leq \Delta_{1} \cdot \Delta_{2} .
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## Bipartite graphs

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## Theorem 18 (Nakprasit, 2008)

Let $G$ be a bipartite graph with maximum degree of partite sets 2 and $\Delta$, then

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■ For $(3, \Delta)$-graphs there is a weaker result: $\chi_{s}^{\prime}(G) \leq 4 \Delta$

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■ Let $M_{i}$ be the set of the edges of same color. Let $G\left(M_{i}\right)$ be a graph induced by $M_{i}$ where every edge from $M_{i}$ is contracted

## Planar graphs

## Theorem 19 (Faudree et al., 1990)

Let $G$ be a planar graph. Then,

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\chi_{s}^{\prime}(G) \leq 4 \Delta(G)+4
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- Altogether we need $4 \chi^{\prime}(G)$ colors


## Planar graphs

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## Planar graphs

■ Forbidding short cycles in planar graphs, gives us some more freedom

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## Conjecture 20 (Hudák et al., 2014)

There exists a constant $C$ such that for every planar graph $G$ with girth $g \geq 5$ it holds

$$
\chi_{s}^{\prime}(G) \leq\left\lceil\frac{2 g(\Delta(G)-1)}{g-1}\right\rceil+C
$$

# Locally irregular edge-coloring 

## Basics

■ A graph $G$ is locally irregular if every two adjacent vertices have distinct degrees.

- An edge-coloring is locally irregular if every color class induces a locally irregular graph.


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■ A graph $G$ is locally irregular if every two adjacent vertices have distinct degrees.

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- Always improper - paths of odd length do not admit such a coloring
- Introduced by Baudon, Bensmail, Przybyło, and Woźniak in 2013 (the paper published in 2015).
- Motivated by the (1-2-3)-Conjecture:

For every graph with no $K_{2}$ component there exists an edge weighting with 1,2 , and 3 such that for every two adjacent vertices the sums on their incident edges are distinct.

## Example: $K_{5}$

■ A test for the audience... How many colors?


Example: $K_{5}$

- Correct!



## Example: $K_{5}$

- And now - add one edge:



## Example: $K_{5}$

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## Decomposable graphs

- A graph is decomposable if it admits a locally irregular edge-coloring (LIE-C).
- The minimum $k$ for which there is a LIE-C of a graph $G$ with $k$ colors is the locally irregular chromatic index of $G, \chi_{\text {irr }}^{\prime}(G)$.


## Decomposable graphs

- A graph is decomposable if it admits a locally irregular edge-coloring (LIE-C).
- The minimum $k$ for which there is a LIE-C of a graph $G$ with $k$ colors is the locally irregular chromatic index of $G, \chi_{\text {irr }}^{\prime}(G)$.
■ Not all graphs are decomposable, e.g. odd-length paths, odd-length cycles.
- A complete characterization was given by Baudon, Bensmail, Przybyło, and Woźniak.


## Decomposable graphs

Define a family of graphs $\mathcal{T}$ recursively:

- The triangle $C_{3}$ belongs to $\mathcal{T}$.
- Every other graph of this family can be constructed by taking an auxiliary graph $F$ which might either be a path of even length or a path of odd length with a triangle glued to one end, then choosing a graph $G \in \mathcal{T}$ containing a triangle with at least one vertex $v$ of degree 2 and finally identifying $v$ with a vertex of degree 1 in $F$.


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## The Conjecture

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Theorem 22 (Baudon et al., 2015)
For every $d$-regular graph $G$, with $d \geq 10^{7}$, it holds $\chi_{\text {irr }}^{\prime}(G) \leq 3$.

Theorem 23 (Przybyło, 2016)
For every graph $G$, with $\delta(G) \geq 10^{10}$, it holds $\chi_{\text {irr }}^{\prime}(G) \leq 3$.

## The upper bound

Bensmail, Merker, and Thomassen established the first constant upper bound using decompositions into bipartite graphs.

## Theorem 24 (Bensmail et al., 2017)

For every decomposable graph $G$, it holds $\chi_{\text {irr }}^{\prime}(G) \leq 328$.
Currently the best:

## Theorem 25 (BL, Przybyło, Soták, 2018+)

For every decomposable graph $G$, it holds $\chi_{\text {irr }}^{\prime}(G) \leq 220$.

## Subcubic graphs

## Theorem 26 (BL, Przybyło, Soták, 2018+)

For every decomposable graph $G$ with $\Delta(G)=3$, it holds $\chi_{\text {irr }}^{\prime}(G) \leq 4$.

## Bipartite graphs

## Theorem 27 (Baudon et al., 2015)

Let $G$ be a regular bipartite graph with minimum degree at least 3 . Then

$$
\chi_{\mathrm{irr}}^{\prime}(G) \leq 2
$$

A decomposable bipartite graph is balanced if all the vertices in one of the two partition parts have even degrees.

## Lemma 28 (Bensmail et al., 2017)

Let $F$ be a balanced forest. Then $F$ admits a LIE-C with at most 2 colors. Moreover, for each vertex $v$ in the partition with no vertex of odd degree, all edges incident to $v$ have the same color.

## Bipartite graphs

## Theorem 29 (BL, Przybyło, Soták, 2018+)

Let $G$ be a (multi)graph not isomorphic to an odd cycle. Then

$$
\chi_{\text {irr }}^{\prime}(\mathcal{S}(G)) \leq 2
$$

Here, $\mathcal{S}(G)$ denotes the full subdivision of $G$, i.e. each edge of $G$ is subdivided once.

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## Question 30

Is every connected balanced graph, which is not a cycle of length $4 k+2$, locally irregularly 2-edge-colorable?

## Bipartite graphs

## Theorem 31 (BL, Przybyło, Soták, 2018+)

Let $G$ be a balanced graph. Then

$$
\chi_{\mathrm{irr}}^{\prime}(G) \leq 4
$$

Theorem 32 (BL, Przybyło, Soták, 2018+)
Let $G$ be a decomposable bipartite graph. Then

$$
\chi_{\mathrm{irr}}^{\prime}(G) \leq 7 .
$$

Moreover, if $G$ has an even number of edges, then the upper bound is 6 .

Ďakujem!

$$
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$$

