

On Incidence Colorings of Graphs

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joint work with

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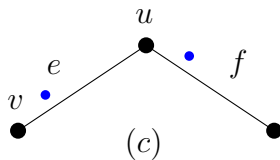
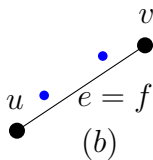
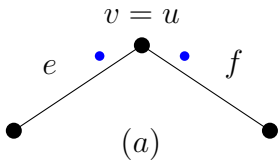
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Incidences

- In a graph G , an **incidence** is a pair (v, e) , where $v \in V(G)$, $e \in E(G)$, and v is incident to e .
- Two incidences (v, e) and (u, f) are **adjacent** if:
 - (a) $v = u$, or
 - (b) $e = f$, or
 - (c) $vu \in \{e, f\}$.



Incidence coloring

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- Related to other colorings, e.g. strong edge-colorings of fully subdivided graphs.
- Example: C_3

Spectrum

- For an incidence coloring c , the **spectrum of a vertex v** , $S_c(v)$, is the set of colors assigned to the incidences with the edges containing v , i.e.

$$S_c(v) = \{c(v, uv), c(u, uv) \mid uv \in E(G)\}.$$

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- Spectrum gives us a simple lower bound:

$$\chi_i(G) \geq \min_c \max_{v \in V(G)} |S_c(v)| \geq \Delta(G) + 1.$$

(k, p) -incidence coloring

- Two types of colors in the spectrum of v :
 - colors of incidences of type (v, uv) - $S_0(v)$;
 - colors of incidences of type (u, uv) - $S_1(v)$.

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- $\chi_i(G) \leq \chi_{i,p}(G)$

$(\Delta + 1)$ -graphs

- A $(\Delta + 1)$ -graph is every graph G with

$$\chi_i(G) = \chi_{i,1}(G) = \Delta(G) + 1.$$

- Complete graphs, trees, outerplanar graphs with $\Delta \geq 7$, planar graphs with $\Delta \geq 14, \dots$
- $\chi_i(G) \leq \chi_{i,1}(G) = \chi(G^2)$

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Theorem 1 (Sun, 2012)

If G is an n -regular graph, then $\chi_i(G) = n + 1$ if and only if $V(G)$ is a disjoint union of $n + 1$ (perfect) dominating sets.

$(\Delta + \ell)$ -graphs

- A $(\Delta + \ell)$ -graph is every graph G with

$$\chi_i(G) \leq \Delta(G) + \ell.$$

- A (k, p) -graph is every graph G with

$$\chi_{i,p}(G) \leq k.$$

$(\Delta + 2)$ -conjecture

Conjecture 2 (Brualdi & Massey, 1993)

For every graph G

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- Guilduli, 1997 - Conjecture is false
- Paley graphs need $\Delta + \Omega(\log \Delta)$

Theorem 3 (Guilduli, 1997)

For every graph G

$$\chi_i(G) \leq \Delta(G) + 20 \log(\Delta(G)) + 84.$$

$(\Delta + 2)$ -conjecture

- Conjecture 2 holds for e.g. **cubic**, partial 2-trees (hence also outerplanar graphs), toroidal grids, planar graphs with girth at least 6 and maximum degree at least 5, complete bipartite graphs,...
- Exists a graph G on **11 vertices** with **maximum degree 6** and

$$\chi_i(G) = 9.$$

- **Open for degrees 4 and 5.**

Theorem 4

For every graph G with maximum degree 4,

$$\chi_i(G) \leq 7.$$

Hypercubes

Theorem 5 (Pai et al., 2014)

For every integers $p, q \geq 1$,

- (i) $\chi_i(Q_n) = n + 1$, if $n = 2^p - 1$;
- (ii) $\chi_i(Q_n) = n + 2$, if $n = 2^p - 2$ and $p \geq 2$, or $n = 2^p + 2^q - 1$, or $n = 2^p + 2^q - 3$ and $p, q \geq 2$.

Hypercubes conjecture

Our motivation:

Conjecture 6 (Pai et al., 2014)

For every $n \geq 1$ such that $n \neq 2^p - 1$ for every integer $p \geq 1$,

$$\chi_i(Q_n) = n + 2.$$

Cartesian products

Observation. Let G and H be arbitrary graphs. Then

$$\chi_i(G \square H) \leq \chi_i(G) + \chi_i(H).$$

- Is it possible that Conjecture 2 holds for Cartesian products?

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- No.

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- Is it possible that Conjecture 2 holds for Cartesian products?
- No.
- Consider a Paley graph P and K_2 ,

$$\chi_i(P \square K_2) = \Delta(P \square K_2) + \Omega(\log(P \square K_2)).$$

Cartesian products

-1 color

Cartesian products (-1 color)

Theorem 7

Let G be a $(\Delta + 1)$ -graph and let H be a subgraph of a regular $(\Delta + 1)$ -graph H' such that

$$\Delta(G) + 1 \geq \Delta(H') - \Delta(H).$$

Then,

$$\chi_i(G \square H) \leq \Delta(G \square H) + 2.$$

Proof of Theorem 7

Proof (1/3).

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- $C = A \cap B = \{t, \dots, \Delta(G)\}$ - overlapping colors between c and d

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- $C = A \cap B = \{t, \dots, \Delta(G)\}$ - overlapping colors between c and d
- $|C| = \Delta(G) - t + 1 = \Delta(H') - \Delta(H)$
- $H \subseteq H'$ and $d = d'|_H \Rightarrow |S_d(v)| = d(v) + 1 \leq \Delta(H) + 1$

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Proof (2/3).

- $M(v) = B \setminus S_d(v)$ - colors from B missing at v

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- \Rightarrow exists an **injective mapping**

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- \Rightarrow exists i.c. f of $G \square H$ with

$$\begin{aligned} |A| + |B| - |C| &= \Delta(G) + 1 + \Delta(H') + 1 - (\Delta(H') - \Delta(H)) \\ &= \Delta(G) + \Delta(H) + 2 \\ &= \Delta(G \square H) + 2 \end{aligned}$$

Proof of Theorem 7

Proof (3/3).

- For every pair of vertices $u \in V(G)$, $v \in V(H)$, and edges $uu' \in E(G)$, $vv' \in E(H)$, we define f :

$$f((u, v), (u, v)(u, v')) = d(v, vv'), \quad \text{and}$$

$$f((u, v), (u, v)(u', v)) = \begin{cases} c(u, uu') & \text{if } c(u, uu') \notin C, \\ g_v(c(u, uu')) & \text{if } c(u, uu') \in C. \end{cases}$$

- We check that f is indeed an i.c. on the board

Hypercubes - revisited

Corollary 8

For every $n \geq 1$,

$$\chi_i(Q_n) = \begin{cases} n + 1 & \text{if } n = 2^m - 1 \text{ for some integer } m \geq 0, \\ n + 2 & \text{otherwise.} \end{cases}$$

Sketch of the proof: $Q_n = Q_{n-k} \square Q_k$, with $k \leq n$, and take $n = 2^m - 1 + k$.

Generalization

(Simplified version)

Theorem 9 (Shiau, Shiau, Wang, 2015)

Let G and H be arbitrary graphs and $s(H)$ the maximum size of a spectrum in H . If $\chi_i(G) \geq \chi_i(H) - s(H)$, then

$$\chi_i(G \square H) \leq \chi_i(G) + s(H).$$

Cartesian products -2 colors

Locally injective homomorphisms

- A **homomorphism** f of G to H is a mapping

$$f : V(G) \rightarrow V(H)$$

such that if $uv \in E(G)$, then $f(u)f(v) \in E(H)$.

- A homomorphism f is **locally injective** if $f(u) \neq f(v)$ for every $v \in V(G)$ and every pair $vu, vw \in E(G)$.
- f is **injective on $N(v)$** , for every $v \in V(G)$
- locally injective homomorphisms **preserve adjacencies of incidences**

Locally injective homomorphisms

Theorem 10 (Duffy, 2015)

Let G and H be simple graphs such that G admits a locally injective homomorphism to H . Then

$$\chi_i(G) \leq \chi_i(H).$$

Proposition 11

A graph G admits a $(k, 1)$ -incidence coloring iff it admits a locally injective homomorphism to K_k .

2-permutable graphs

- K_{2n}^- is a complete graph on $2n$ vertices without a perfect matching
- A connected $2d$ -regular graph G is **2-permutable** if it admits a locally injective homomorphism to K_{2d+2}^- .

2-permutable graphs

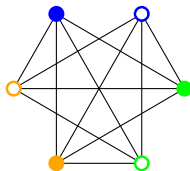
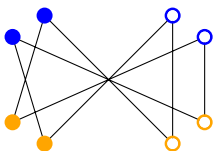
- K_{2n}^- is a complete graph on $2n$ vertices without a perfect matching
- A connected $2d$ -regular graph G is **2-permutable** if it admits a locally injective homomorphism to K_{2d+2}^- .
- So:
 - G is **$(2d + 2)$ -partite** (with partition sets P_1, \dots, P_{2d+2});
 - For every i , $1 \leq i \leq 2d + 2$, exists \bar{i} such that there are **no edges between P_i and $P_{\bar{i}}$** ;
 - Every $v \in P_i$ has **at most one neighbor in P_j , $j \notin \{i, \bar{i}\}$** .
 - Every 2-permutable graph is a **$(\Delta + 2, 1)$ -graph**.

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 - Every $v \in P_i$ has **at most one neighbor in P_j , $j \notin \{i, \bar{i}\}$** .
 - Every 2-permutable graph is a **$(\Delta + 2, 1)$ -graph**.
- There exist $(\Delta + 2, 1)$ -graphs which are not 2-permutable, e.g. 7-cycle.

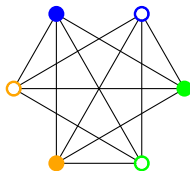
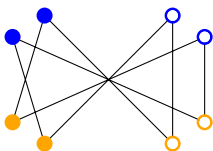
2-permutable graphs

Examples: C_{4n} and K_{2n}^-



2-permutable graphs

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- Among 1544 4-regular graphs of order 12, there are 13 2-permutable graphs.

Prisms over 2-permutable graphs

Theorem 12

Let G be a 2-permutable graph. Then

$$\chi_i(G \square K_2) = \Delta(G \square K_2) + 1 (= \Delta(G) + 2).$$

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- The prism over the Dodecahedron is a $(\Delta + 1)$ -graph, while the Dodecahedron is not 2-permutable (it is cubic).
- The inverse holds for cycles.

Sub-2-permutable graphs

- A (non-regular) graph G is **sub-2-permutable** if it admits a locally injective homomorphism to $K_{\Delta(G)+2}^-$.

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Corollary 13

Let G be a sub-2-permutable graph. Then

$$\chi_i(G \square K_2) = \Delta(G \square K_2) + 1.$$

2-adjustable graphs

- An incidence coloring of a graph G is **adjustable** if there exists a pair of colors x and y such that there is no vertex $v \in V(G)$ with $x, y \in S_0(v)$.
- x and y are called **free colors**.

2-adjustable graphs

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- A graph G is **2-adjustable** if it admits an adjustable $(\Delta(G) + 2)$ -incidence coloring.

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- Example: C_5

2-adjustable graphs - Examples

- All $(\Delta + 1)$ -graphs (the color $\Delta(G) + 2$ is never used).

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- All $(\Delta + 1)$ -graphs together with a matching (two same colors can be put on a matching; they are free)

2-adjustable graphs - Examples

- All $(\Delta + 1)$ -graphs (the color $\Delta(G) + 2$ is never used).
- All $(\Delta + 1)$ -graphs together with a matching (two same colors can be put on a matching; they are free)
- Cycles, complete bipartite graphs, prisms over C_{6n}

2-adjustable graphs - Examples

- By \mathring{K}_n we denote the complete graph of order n with a loop at one vertex.

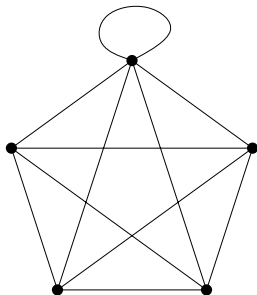


Figure: A \mathring{K}_5 .

2-adjustable graphs - Examples

Proposition 14

If a graph G admits a locally injective homomorphism to $K_{\Delta(G)+1}$, then G is 2-adjustable.

2-adjustable graphs - Examples

Proposition 14

If a graph G admits a locally injective homomorphism to $\mathring{K}_{\Delta(G)+1}$, then G is 2-adjustable.

- The inverse statement is not true in general.
- C_5 is 2-adjustable, but does not admit a locally injective homomorphism to \mathring{K}_3 .

Cartesian products with -2 colors

Theorem 15

Let G be a sub-2-permutable graph and let H be a 2-adjustable graph. Then

$$\chi_i(G \square H) \leq \Delta(G \square H) + 2.$$

Open problems

Conjecture 16

Let G be a $(\Delta + 1)$ -graph and H be a $(\Delta + 2)$ -graph. Then,

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Let G be a $(\Delta + 1)$ -graph and H be a $(\Delta + 2)$ -graph. Then,

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Question 17

Do there exist graphs G and H with $\chi_i(G) = \Delta(G) + 2$ and $\chi_i(H) = \Delta(H) + 2$ such that $\chi_i(G \square H) > \Delta(G \square H) + 2$?

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Question 18

When is the Cartesian product of two $(\Delta + 1)$ -graphs also a $(\Delta + 1)$ -graph.

Hvala!