# On Incidence Colorings of Graphs 

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joint work with

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## Graph Colorings

## Vertices

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Regions of every map can be colored with at most four colors such that every pair of regions with a common border is colored differently.

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## Edges

■ Vizing, 1964
The edges of every simple graph $G$ can be colored with at most $\Delta(G)+1$ colors such that incident edges are colored differently.

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■ Applicable (up to some level) in practice for solving optimization problems

- Most of coloring problems are NP-complete


## Incidences

■ In a graph $G$, an incidence is a pair $(v, e)$, where $v \in V(G)$, $e \in E(G)$, and $v$ is incident to $e$.

- Two incidences $(v, e)$ and $(u, f)$ are adjacent if:

$$
\begin{aligned}
& \text { (a) } v=u \text {, or } \\
& \text { (b) } e=f, \text { or } \\
& \text { (c) } v u \in\{e, f\} \text {. }
\end{aligned}
$$



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■ Defined by Brualdi and Massey in 1993.
- Related to other colorings, e.g. strong edge-coloring of fully subdivided graphs.

Example: $C_{3}$



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## Spectrum

- For an incidence coloring $c$, the spectrum of a vertex $v$, $S_{c}(v)$, is the set of colors assigned to the incidences with the edges containing $v$, i.e.

$$
S_{c}(v)=\{c(v, u v), c(u, u v) \mid u v \in E(G)\} .
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- $\left|S_{c}^{0}(v)\right|=d(v)$ and $\left|S_{c}^{1}(v)\right| \geq 1$


## Spectrum

- $S_{c}^{0}(v)=\{c(v, u v) \mid u v \in E(G)\}$
- $S_{c}^{1}(v)=\{c(u, u v) \mid u v \in E(G)\}$
- $\left|S_{c}^{0}(v)\right|=d(v)$ and $\left|S_{c}^{1}(v)\right| \geq 1$
- Spectrum gives a simple lower bound:

$$
\chi_{i}(G) \geq \min _{c} \max _{v \in V(G)}\left|S_{c}(v)\right| \geq \Delta(G)+1
$$

## $(\Delta+1)$-graphs

- A $(\Delta+1)$-graph is every graph $G$ with

$$
\chi_{i}(G)=\Delta(G)+1
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- Complete graphs, trees, outerplanar graphs with $\Delta \geq 7$, planar graphs with $\Delta \geq 14, \ldots$


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## Theorem 1 (Sun, 2012)

If $G$ is an n-regular graph, then $\chi_{i}(G)=n+1$ if and only if $V(G)$ is a disjoint union of $n+1$ (perfect) dominating sets.

## $(\Delta+2)$-conjecture

Conjecture 2 (Brualdi \& Massey, 1993)
For every graph G

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## Theorem 3 (Guilduli, 1997)

For every graph G

$$
\chi_{i}(G) \leq \Delta(G)+20 \log (\Delta(G))+84
$$

## $(\Delta+2)$-conjecture

- Conjecture 2 holds for e.g.
- subcubic graphs,
- partial 2-trees (hence also outerplanar graphs),
- toroidal grids,
- planar graphs with girth at least 6 and maximum degree at least 5,
- complete bipartite graphs,


## $(\Delta+2)$-conjecture

The graph $G$ of smallest order being a counter example (Clark \& Dunning, 1997):
6 -regular, 11 vertices, $\gamma(G)=3, \chi_{i}(G)=9$


## $(\Delta+2)$-conjecture

## Theorem 4 (Maydanskiy, 2005)

$$
\chi_{i}(G) \geq \frac{2|E(G)|}{|V(G)|-\gamma(G)}
$$

- So far, the only known graphs being counter-examples to the conjecture are the ones having high domination number
- Open for $\Delta(G) \in\{4,5\}$.
- Strong edge-coloring result gives $\chi_{i}(G) \leq 2 \Delta(G)$ (Nakprasit, 2008)


## Subquartic graphs

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For every graph $G$ with maximum degree 4,

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- 4-regular graphs on at most 14 vertices are $(\Delta+2)$-graphs;

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■ [many 4-regular graphs on 15 vertices are ( $\Delta+2$ )-graphs] :)

## Question 6

Is it true that 6 (resp. 7) colors suffice for incidence coloring of graphs with maximum degree 4 (resp. 5)?

## Hypercubes

## Theorem 7 (Pai et al., 2014)

For every integers $p, q \geq 1$,
(i) $\chi_{i}\left(Q_{n}\right)=n+1$, if $n=2^{p}-1$;
(ii) $\chi_{i}\left(Q_{n}\right)=n+2$, if $n=2^{p}-2$ and $p \geq 2$, or $n=2^{p}+2^{q}-1$, or $n=2^{p}+2^{q}-3$ and $p, q \geq 2$.

## Hypercubes

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## Our motivation:

## Conjecture 8 (Pai et al., 2014)

For every $n \geq 1$ such that $n \neq 2^{p}-1$ for every integer $p \geq 1$,

$$
\chi_{i}\left(Q_{n}\right)=n+2 .
$$

## Cartesian products

Observation. Let $G$ and $H$ be arbitrary graphs. Then

$$
\chi_{i}(G \square H) \leq \chi_{i}(G)+\chi_{i}(H)
$$

- Is it possible that Conjecture 2 holds for Cartesian products?


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## Cartesian products

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- Is it possible that Conjecture 2 holds for Cartesian products?
- No.
- Consider a Paley graph $P$ and $K_{2}$,

$$
\chi_{i}\left(P \square K_{2}\right)=\Delta\left(P \square K_{2}\right)+\Omega\left(\log \left(P \square K_{2}\right)\right) .
$$

## Cartesian products -1 color

## Cartesian products ( -1 color)

## Theorem 9 (Gregor, •, Soták, 2016)

Let $G$ be a $(\Delta+1)$-graph and let $H$ be a subgraph of a regular $(\Delta+1)$-graph $H^{\prime}$ such that

$$
\Delta(G)+1 \geq \Delta\left(H^{\prime}\right)-\Delta(H)
$$

Then,

$$
\chi_{i}(G \square H) \leq \Delta(G \square H)+2 .
$$

## Hypercubes - revisited

Corollary 10 (Gregor, •, Soták, 2016)
For every $n \geq 1$,

$$
\chi_{i}\left(Q_{n}\right)= \begin{cases}n+1 & \text { if } n=2^{m}-1 \text { for some integer } m \geq 0 \\ n+2 & \text { otherwise }\end{cases}
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The conjecture has also been solved independently by Shiau, Shiau, Wang, 2015

## Cartesian products <br> -2 colors

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Not today :(

## Open problems

## Conjecture 11

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## Open problems

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## Question 12

Do there exist graphs $G$ and $H$ with $\chi_{i}(G)=\Delta(G)+2$ and $\chi_{i}(H)=\Delta(H)+2$ such that $\chi_{i}(G \square H)>\Delta(G \square H)+2$ ?

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## Question 13

When is the Cartesian product of two $(\Delta+1)$-graphs also a ( $\Delta+1$ )-graph.

## Thank you for your attention!

## Locally injective homomorphisms

- A homomorphism $f$ of $G$ to $H$ is a mapping

$$
f: V(G) \rightarrow V(H)
$$

such that if $u v \in E(G)$, then $f(u) f(v) \in E(H)$.

- A homomorphism $f$ is locally injective if $f(u) \neq f(v)$ for every $v \in V(G)$ and every pair $v u, v w \in E(G)$.
- $f$ is injective on $N(v)$, for every $v \in V(G)$
- locally injective homomorphisms preserve adjacencies of incidences


## Locally injective homomorphisms

## Theorem 14 (Duffy, 2015)

Let $G$ and $H$ be simple graphs such that $G$ admits a locally injective homomorphism to $H$. Then

$$
\chi_{i}(G) \leq \chi_{i}(H)
$$

## Proposition 15

A graph $G$ admits a ( $k, 1$ )-incidence coloring iff it admits a locally injective homomorphism to $K_{k}$.

## 2-permutable graphs

- $K_{2 n}^{-}$is a complete graph on $2 n$ vertices without a perfect matching
- A connected $2 d$-regular graph $G$ is 2 -permutable if it admits a locally injective homomorphism to $K_{2 d+2}^{-}$.


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- A connected $2 d$-regular graph $G$ is 2 -permutable if it admits a locally injective homomorphism to $K_{2 d+2}^{-}$.
- So:

■ $G$ is $(2 d+2)$-partite (with partition sets $\left.P_{1}, \ldots, P_{2 d+2}\right)$;
■ For every $i, 1 \leq i \leq 2 d+2$, exists $\bar{i}$ such that there are no edges between $P_{i}$ and $P_{\bar{i}}$;
■ Every $v \in P_{i}$ has at most one neighbor in $P_{j}, j \notin\{i, \bar{i}\}$.
■ Every 2 -permutable graph is a $(\Delta+2,1)$-graph.

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- Every $v \in P_{i}$ has at most one neighbor in $P_{j}, j \notin\{i, \bar{i}\}$.
- Every 2 -permutable graph is a $(\Delta+2,1)$-graph.
- There exist $(\Delta+2,1)$-graphs which are not 2-permutable, e.g. 7-cycle.


## 2-permutable graphs

Examples: $C_{4 n}$ and $K_{2 n}^{-}$


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- Among 1544 4-regular graphs of order 12, there are 13 2-permutable graphs.


## Prisms over 2-permutable graphs

## Theorem 16

Let $G$ be a 2-permutable graph. Then

$$
\chi_{i}\left(G \square K_{2}\right)=\Delta\left(G \square K_{2}\right)+1(=\Delta(G)+2)
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- The prism over the Dodecahedron is a ( $\Delta+1$ )-graph, while the Dodecahedron is not 2-permutable (it is cubic).
- The inverse holds for cycles.


## Sub-2-permutable graphs

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## Corollary 17

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## 2-adjustable graphs

- An incidence coloring of a graph $G$ is adjustable if there exists a pair of colors $x$ and $y$ such that there is no vertex $v \in V(G)$ with $x, y \in S_{0}(v)$.
- $x$ and $y$ are called free colors.


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- Example: $C_{5}$


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■ All $(\Delta+1)$-graphs (the color $\Delta(G)+2$ is never used).

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- All $(\Delta+1)$-graphs together with a matching (two same colors can be put on a matching; they are free)
■ Cycles, complete bipartite graphs, prisms over $C_{6 n}$


## 2-adjustable graphs - Examples

- By $\stackrel{\circ}{K}_{n}$ we denote the complete graph of order $n$ with a loop at one vertex.


Figure: A $\stackrel{\circ}{K}_{5}$.

## 2-adjustable graphs - Examples

## Proposition 18

If a graph $G$ admits a locally injective homomorphism to ${\stackrel{\circ}{K_{\Delta(G)+1}}}$, then $G$ is 2-adjustable.

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- The inverse statement is not true in general.
- $C_{5}$ is 2 -adjustable, but does not admit a locally injective homomorphism to $\check{K}_{3}$.


## Cartesian products with -2 colors

## Theorem 19

Let $G$ be a sub-2-permutable graph and let $H$ be a 2-adjustable graph. Then

$$
\chi_{i}(G \square H) \leq \Delta(G \square H)+2 .
$$

## Open problems

## Conjecture 20

Let $G$ be a $(\Delta+1)$-graph and $H$ be a $(\Delta+2)$-graph. Then,

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